

A Coinductive Version of Milner's Proof System for Regular Expressions Modulo Bisimilarity

Clemens Grabmayer

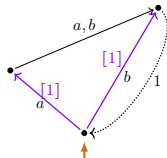
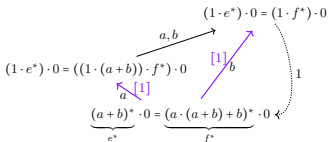
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Process semantics $\llbracket \cdot \rrbracket_P$ of regular (star) expr's *(Milner, 1984)*

$0 \xrightarrow{P}$ deadlock δ , no termination

$1 \xrightarrow{P}$ empty process ϵ , then terminate

$a \xrightarrow{P}$ atomic action a , then terminate

$e + f \xrightarrow{P}$ alternative composition of $\llbracket e \rrbracket_P$ and $\llbracket f \rrbracket_P$

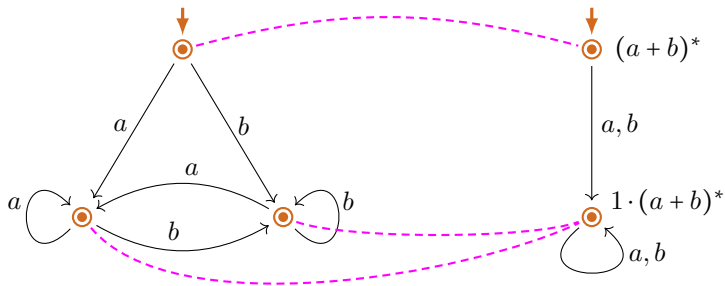
$e \cdot f \xrightarrow{P}$ sequential composition of $\llbracket e \rrbracket_P$ and $\llbracket f \rrbracket_P$

$e^* \xrightarrow{P}$ unbounded iteration of $\llbracket e \rrbracket_P$, option to terminate

$\llbracket e \rrbracket_P := [\mathbf{P}(e)]_{\Leftrightarrow}$ (bisimilarity equivalence class of process $\mathbf{P}(e)$)

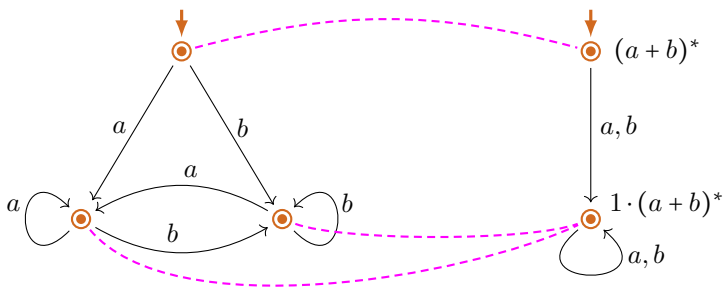
$:= [\mathcal{C}(e)]_{\Leftrightarrow}$ (bisimilarity equivalence class of chart $\mathcal{C}(e)$)

Chart interpretation (example) (via TSS or *Antimirov's partial deriv's*)



$$\begin{aligned} \mathcal{C}((a^* \cdot b^*)^*) & \quad \Longleftrightarrow \quad \mathcal{C}((a+b)^*) \\ \llbracket (a^* \cdot b^*)^* \rrbracket_{\mathcal{P}} & \quad = \quad \llbracket (a+b)^* \rrbracket_{\mathcal{P}} \end{aligned}$$

Chart interpretation (example) (via TSS or *Antimirov's partial deriv's*)



$$\begin{array}{ccc}
 \mathcal{C}((a^* \cdot b^*)^*) & \stackrel{\leftrightarrow}{=} & \mathcal{C}((a+b)^*) \\
 (a^* \cdot b^*)^* & \stackrel{=}{=} \llbracket \cdot \rrbracket_P & (a+b)^*
 \end{array}$$

Milner's proof system $Mil = Mil^- + RSP^*$

Axioms:

$$(assoc(+)) \quad (e + f) + g = e + (f + g)$$

$$(neutr(+)) \quad e + 0 = e$$

$$(comm(+)) \quad e + f = f + e$$

$$(idempot(+)) \quad e + e = e$$

$$(assoc(\cdot)) \quad (e \cdot f) \cdot g = e \cdot (f \cdot g)$$

$$(r-distr(+, \cdot)) \quad (e + f) \cdot g = e \cdot g + f \cdot g$$

$$(id_l(\cdot)) \quad 1 \cdot e = e$$

$$(id_r(\cdot)) \quad e \cdot 1 = e$$

$$(deadlock) \quad 0 \cdot e = 0$$

$$(rec(*)) \quad e^* = 1 + e \cdot e^*$$

$$(trm-body(*)) \quad e^* = (1 + e)^*$$

Inference rules: equational logic plus

$$\frac{e = f \cdot e + g}{e = f^* \cdot g} RSP^* \quad (\text{if } f \downarrow)$$

Milner's axiomatization question

Question (Milner, 1984)

Is Milner's system $\text{Mil} = \text{Mil}^- + \text{RSP}^*$ **complete** for bisimilarity of process interpretations of regular expressions?

$$\forall e, f \text{ reg. expr's } \left(\vdash_{\text{Mil}} e = f \begin{array}{c} \xleftarrow{\text{complete?}} \\ \xrightarrow{\text{sound}} \end{array} e = \llbracket \cdot \rrbracket_P f \right) ?$$

"Yes" for restrictions to subclasses: *Zantema/Fokkink (1994), Fokkink (1996), Corradini, De Nicola, Labella (2002), G/Fokkink (2020)*.

Proposition (G, CMCS 2006)

The system $\text{Mil}^- + \text{USP}$, where:

USP: **unique solvability** of guarded, linear systems of equations, is (sound and) **complete**.

Question (investigated here)

How can the derivational power be characterized that the **fixed-point rule RSP*** adds to the purely equational part Mil^- of Milner's system?

Answer developed

We use:

- ▶ the loop existence and elimination property (LEE) of charts
 - ▶ implies expressibility by a star expression
 - ▶ led to completeness result for 1-free star expressions (G/Fokkink, 2020)

We introduce:

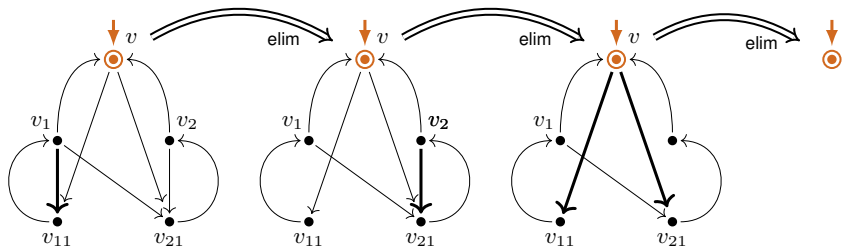
- ▶ a coinductive version $\text{cMil} = (\text{Mil}^- + \text{LCoProof})$ of $\text{Mil} = (\text{Mil}^- + \text{RSP}^*)$ based on LEE-witnessed coinductive proofs over Mil^- .

We construct / obtain:

- ▶ a proof transformation: $\text{Mil} \mapsto \text{cMil}$, (RSP^* inst's \mapsto LCoProof inst's),
- ▶ a proof transformation: $\text{Mil} \longleftarrow \text{cMil}$, (bottom-up extraction procedure),
- ▶ theorem equivalence $\text{Mil} \sim \text{cMil}$:

$$\vdash_{\text{Mil}} e = f \iff \vdash_{\text{cMil}} e = f.$$

LEE

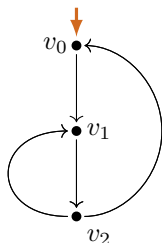


Loop charts (interpretations of innermost iterations in 1-free expressions)

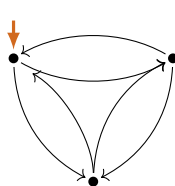
Definition

A chart is a **loop chart** if:

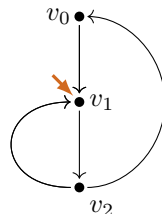
- (L1) There is an infinite path from the **start vertex**.
- (L2) Every infinite path from the **start vertex** returns to it.
- (L3) Immediate termination is **only** possible at the **start vertex**.



(L1), ~~(L2)~~



(L1), (L2), ~~(L3)~~

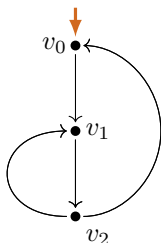


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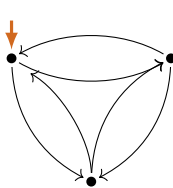
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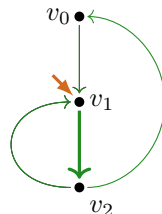
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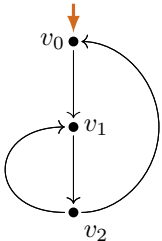


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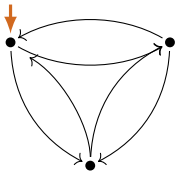
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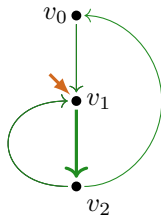
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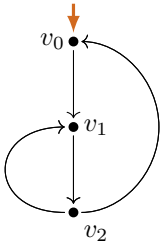
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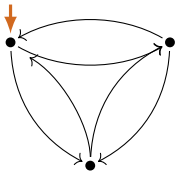
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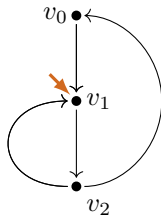
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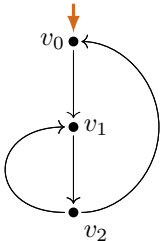
loop chart

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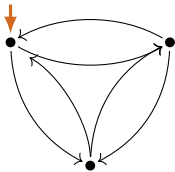
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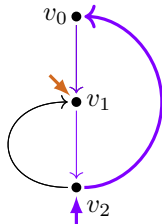
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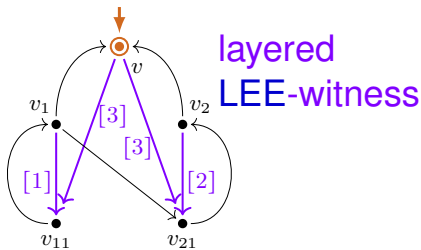
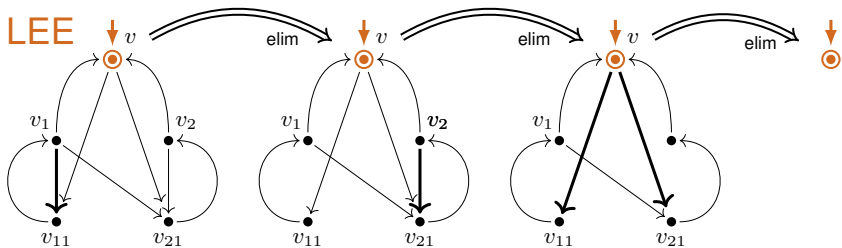


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loop subchart

LEE, and LLEE-witness

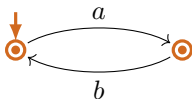


LEE

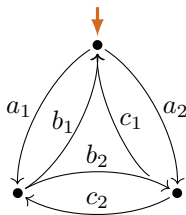
Definition

A chart \mathcal{C} satisfies **LEE** (*loop existence and elimination*) if:

$$\exists \mathcal{C}_0 \left(\mathcal{C} \xRightarrow{*}_{\text{elim}} \mathcal{C}_0 \not\xrightarrow{\text{elim}} \right. \\ \left. \wedge \mathcal{C}_0 \text{ permits no infinite path} \right).$$

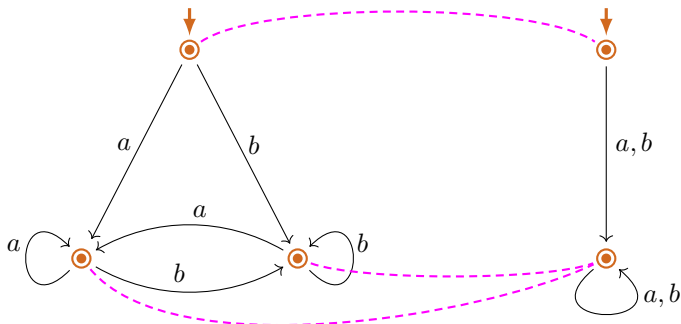


LEE



LEE

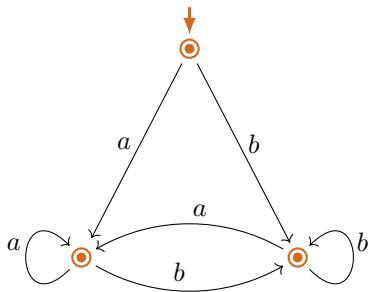
Chart interpretation (example)



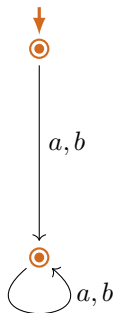
$$\mathcal{C}((a^* \cdot b^*)^*) \quad \Leftrightarrow \quad \mathcal{C}((a + b)^*)$$

$$\llbracket (a^* \cdot b^*)^* \rrbracket_{\mathcal{P}} \quad = \quad \llbracket (a + b)^* \rrbracket_{\mathcal{P}}$$

Chart interpretation (example)

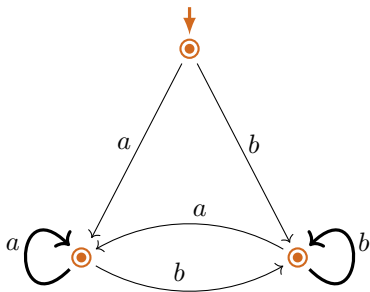


$$\mathcal{C}((a^* \cdot b^*)^*)$$

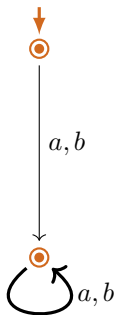


$$\mathcal{C}((a + b)^*)$$

Chart interpretation (example)



$$\mathcal{C}((a^* \cdot b^*)^*)$$



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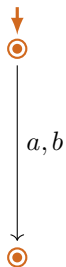
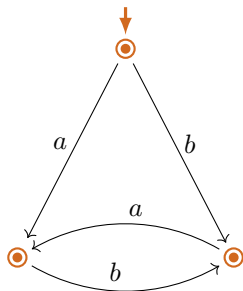
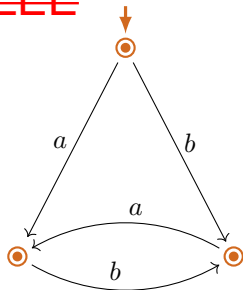


Chart interpretation (example)

LEE



LEE

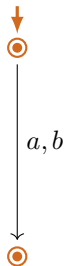
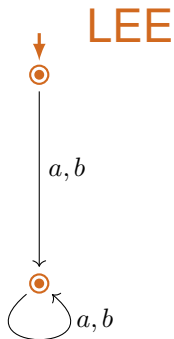
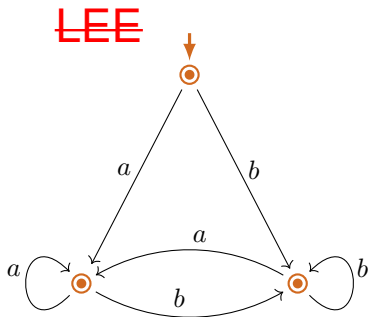
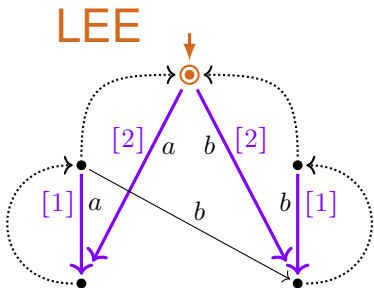


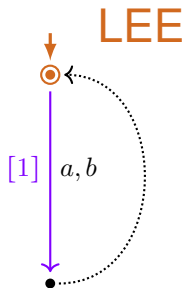
Chart interpretation (example)



LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

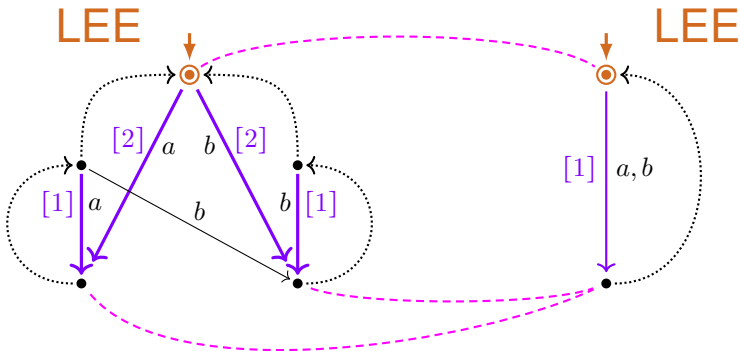


$$\underline{C}((a^* \cdot b^*)^*)$$



$$\underline{C}((a + b)^*)$$

LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

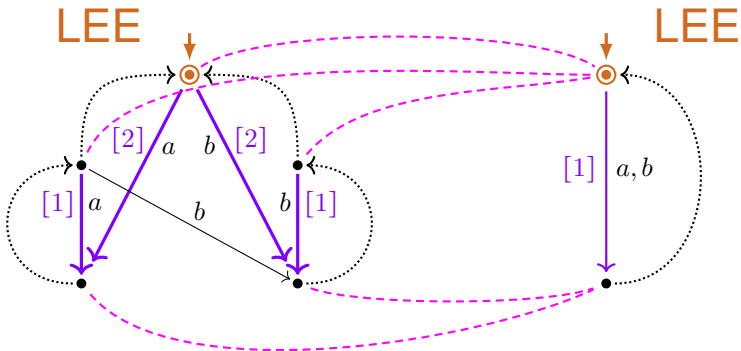


$$\underline{\mathcal{C}}((a^* \cdot b^*)^*)$$

$$\longleftrightarrow^{(1)}$$

$$\underline{\mathcal{C}}((a + b)^*)$$

LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

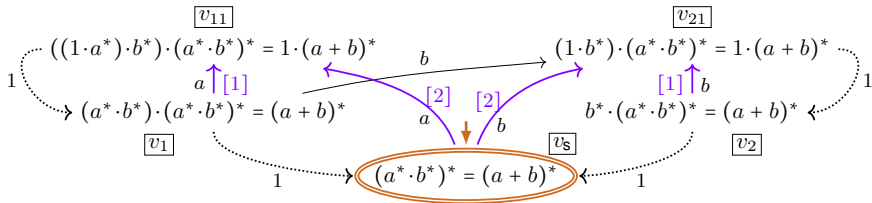


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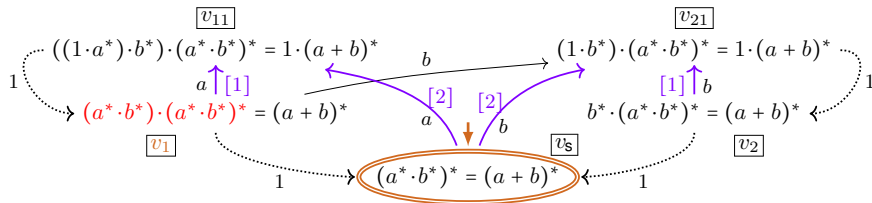
$$\underline{\mathcal{C}}((a + b)^*)$$

LEE-witnessed coinductive proof over Mil⁻



Right- and left-hand sides are Mil⁻-provable solutions in every vertex.

LEE-witnessed coinductive proof over Mil⁻

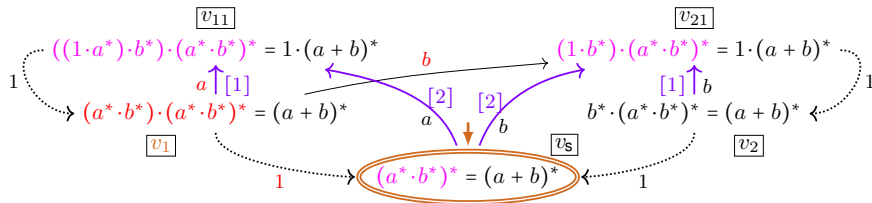


Right- and left-hand sides are Mil⁻-provable solutions in every vertex.

E.g. in v_1 :

$$(a^* \cdot b^*) \cdot (a^* \cdot b^*)^* =_{\text{Mil}^-}$$

LEE-witnessed coinductive proof over Mil⁻

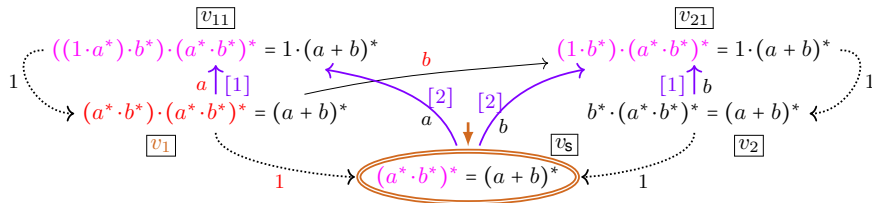


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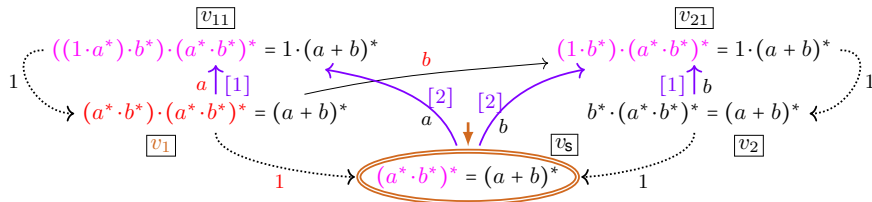
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E.g. in v_1 :

$$(a^* \cdot b^*) \cdot (a^* \cdot b^*)^* =_{\text{Mil}^-}$$

$$=_{\text{Mil}^-} 1 \cdot (a^* \cdot b^*)^* + a \cdot (((1 \cdot a^*) \cdot b^*) \cdot (a^* \cdot b^*)^*) + b \cdot ((1 \cdot b^*) \cdot (a^* \cdot b^*)^*)$$

LEE-witnessed coinductive proof over Mil⁻

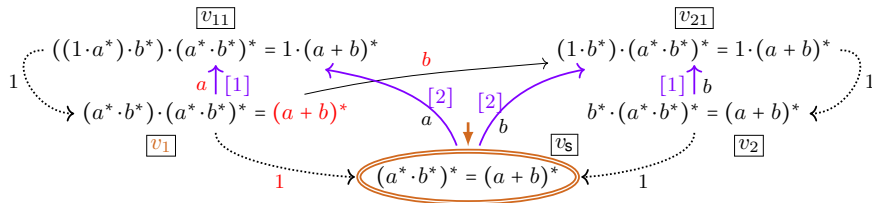


Right- and left-hand sides are **Mil⁻-provable solutions** in every vertex.

E.g. in v_1 :

$$\begin{aligned}
 (a^* \cdot b^*) \cdot (a^* \cdot b^*)^* &=_{\text{Mil}^-} ((1 + a \cdot a^*) \cdot (1 + b \cdot b^*)) \cdot (a^* \cdot b^*)^* \\
 &=_{\text{Mil}^-} (1 \cdot 1 + a \cdot a^* \cdot 1 + 1 \cdot b \cdot b^* + a \cdot a^* \cdot b \cdot b^*) \cdot (a^* \cdot b^*)^* \\
 &=_{\text{Mil}^-} (1 + a \cdot a^* + a \cdot a^* \cdot b \cdot b^* + b \cdot b^*) \cdot (a^* \cdot b^*)^* \\
 &=_{\text{Mil}^-} (1 + a \cdot a^* \cdot (1 + b \cdot b^*) + b \cdot b^*) \cdot (a^* \cdot b^*)^* \\
 &=_{\text{Mil}^-} (1 + a \cdot a^* \cdot b^* + b \cdot b^*) \cdot (a^* \cdot b^*)^* \\
 &=_{\text{Mil}^-} 1 \cdot (a^* \cdot b^*)^* + a \cdot (((1 \cdot a^*) \cdot b^*) \cdot (a^* \cdot b^*)^*) + b \cdot ((1 \cdot b^*) \cdot (a^* \cdot b^*)^*)
 \end{aligned}$$

LEE-witnessed coinductive proof over Mil⁻

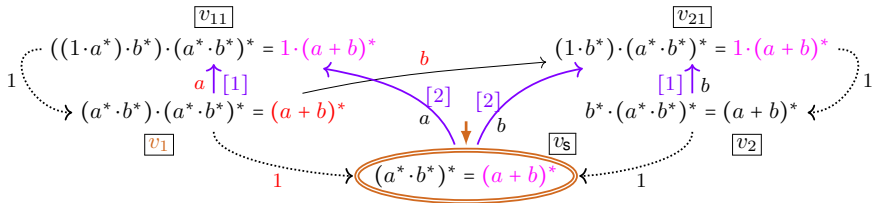


Right- and left-hand sides are Mil⁻-provable solutions in every vertex.

E.g. in v_1 :

$$(a + b)^* =_{\text{Mil}^-}$$

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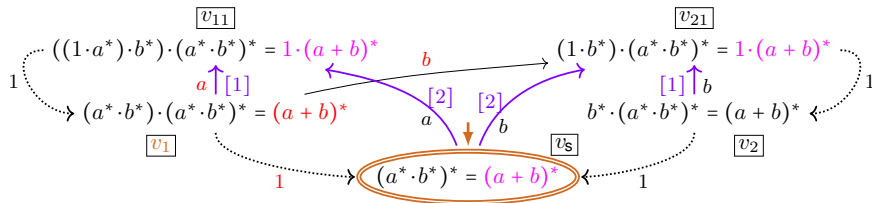


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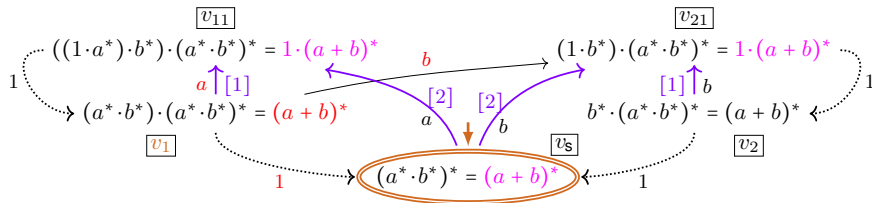
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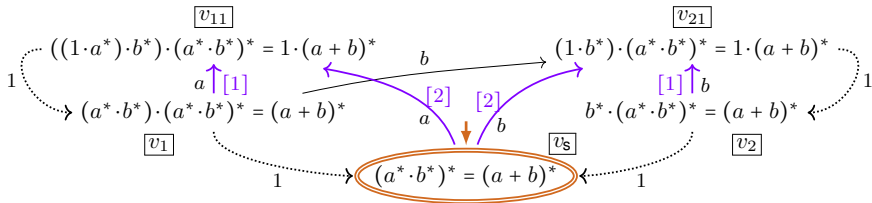


Right- and left-hand sides are Mil⁻-provable solutions in every vertex.

E.g. in v_1 :

$$\begin{aligned}
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 &=_{\text{Mil}^-} 1 + 1 + (a + b) \cdot (a + b)^* + a \cdot (a + b)^* + b \cdot (a + b)^* \\
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 &=_{\text{Mil}^-} 1 \cdot (a + b)^* + a \cdot (1 \cdot (a + b)^*) + b \cdot (1 \cdot (a + b)^*)
 \end{aligned}$$

LEE-witnessed coinductive proof over Mil⁻



Right- and left-hand sides are Mil⁻-provable solutions in every vertex.

Coinductive proof systems

Rule scheme for combining LEE-witnessed coinductive proofs:

$$\frac{g_1 = h_1 \quad \dots \quad g_n = h_n \quad \mathcal{LCP}(e = f)}{e = f} \text{LCoProof}_n$$

- ▶ $\mathcal{LCP}(e = f)$ is a LEE-witnessed coinductive proof of $e = f$ over $\text{Mil}^- + \{g_1 = h_1, \dots, g_n = h_n\}$.

We define the proof systems:

$\text{CLC} := \text{rules } \{\text{LCoProof}_n\}_{n \in \mathbb{N}}$

$\text{cMil} := \text{Mil}^- + \{\text{LCoProof}_n\}_{n \in \mathbb{N}}$

Lemma

$\text{CLC} \sim \text{cMil}$

Coinductive proof systems

Rule scheme for **combining coinductive proofs**:

$$\frac{g_1 = h_1 \quad \dots \quad g_n = h_n \quad \mathcal{CP}(e = f)}{e = f} \text{CoProof}_n$$

- $\mathcal{CP}(e = f)$ is a coinductive proof of $e = f$
over $\text{Mil}^- + \{g_1 = h_1, \dots, g_n = h_n\}$.

We define the proof systems:

$\text{CLC} := \text{rules } \{\text{LCoProof}_n\}_{n \in \mathbb{N}}$

$\text{CC} := \text{rules } \{\text{CoProof}_n\}_{n \in \mathbb{N}}$

$\text{cMil} := \text{Mil}^- + \{\text{LCoProof}_n\}_{n \in \mathbb{N}}$

$\overline{\text{cMil}} := \text{Mil}^- + \{\text{CoProof}_n\}_{n \in \mathbb{N}}$

Lemma

$\text{CLC} \sim \text{cMil}$

Lemma

- (i) $\text{CC} \sim \overline{\text{cMil}}$.
- (ii) $\overline{\text{cMil}} \sim \text{Mil}^- + \text{USP}$,
- (iii) CC and $\overline{\text{cMil}}$
are complete for $=_{\llbracket \cdot \rrbracket_P}$.

Coinductive proof systems

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We define the proof systems:

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$\text{cMil} := \text{Mil}^- + \{\text{LCoProof}_n\}_{n \in \mathbb{N}}$

$\overline{\text{cMil}} := \text{Mil}^- + \{\text{CoProof}_n\}_{n \in \mathbb{N}}$

Lemma

$\text{CLC} \sim \text{cMil}$

Consequence

$\text{cMil} \sim \text{CLC} \approx \text{CC} \sim \overline{\text{cMil}}$.

Lemma

- (i) $\text{CC} \sim \overline{\text{cMil}}$.
- (ii) $\overline{\text{cMil}} \sim \text{Mil}^- + \text{USP}$,
- (iii) CC and $\overline{\text{cMil}}$
are complete for $=_{[\cdot]_P}$.

Proof transformation $\text{Mil} \mapsto \text{cMil}$

$$\begin{array}{c}
 \cdot \\
 \frac{e = f \cdot e + g}{e = f^* \cdot g} \text{RSP}^* \\
 \Longrightarrow \frac{e = f \cdot e + g \quad \mathcal{LCP}_{\text{Mil}^- + \{e = f \cdot e + g\}}(e = f^* \cdot g)}{e = f^* \cdot g} \text{LCoProof}_1
 \end{array}$$

Proof transformation $\text{Mil} \mapsto \text{cMil}$

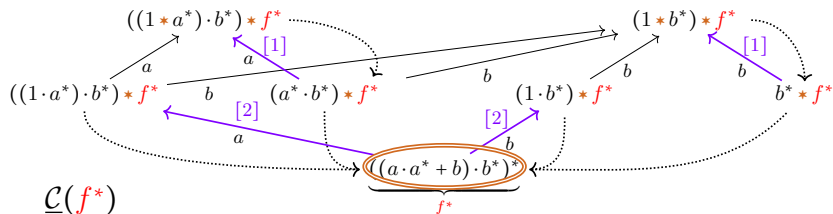
$$\frac{\overbrace{(a+b)^*}^e = \overbrace{((a \cdot a^* + b) \cdot b^*)}^f \cdot \overbrace{(a+b)^*}^e + \overbrace{1}^g}}{\overbrace{(a+b)^*}^e = \underbrace{((a \cdot a^* + b) \cdot b^*)^*}_{f^*} \cdot \underbrace{1}_g} \text{RSP}^*$$

From fixed-point rule instances to coinductive proofs

$$\frac{\overbrace{(a+b)^*}^e = \overbrace{((a \cdot a^* + b) \cdot b^*)}^f \cdot \overbrace{(a+b)^*}^e + \overbrace{1}^g}}{\overbrace{(a+b)^*}^e = \underbrace{((a \cdot a^* + b) \cdot b^*)^*}_{f^*} \cdot \underbrace{1}_g} \text{RSP}^*$$

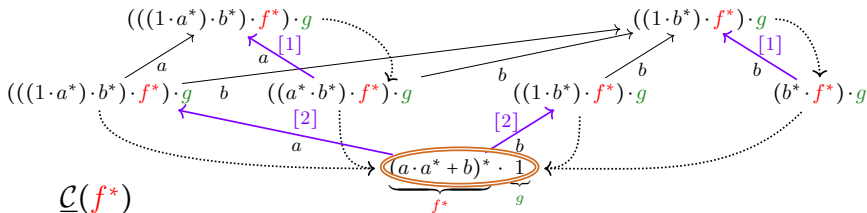
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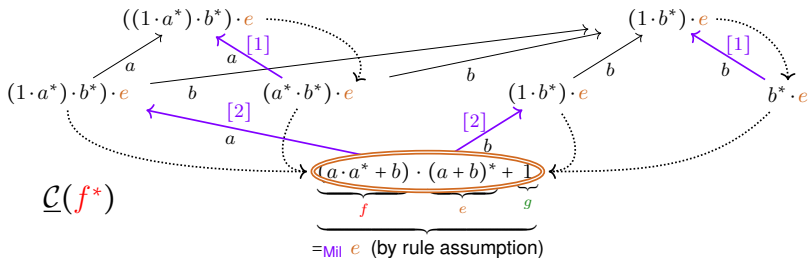
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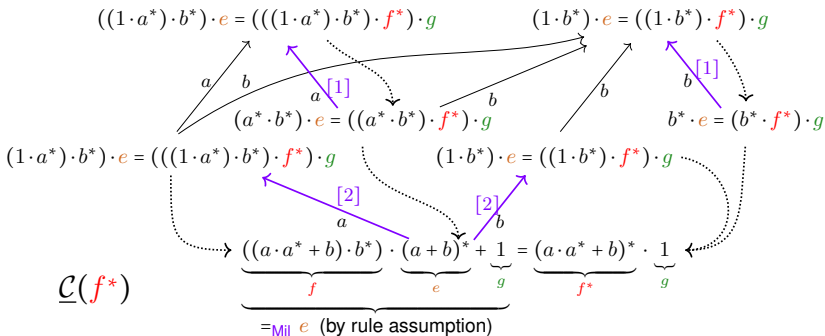
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From fixed-point rule instances to coinductive proofs

$$\frac{\overbrace{(a+b)^*}^e = \overbrace{((a \cdot a^* + b) \cdot b^*)}^f \cdot \overbrace{(a+b)^*}^e + \overbrace{1}^g}}{\overbrace{(a+b)^*}^e = \underbrace{((a \cdot a^* + b) \cdot b^*)^*}_{f^*} \cdot \underbrace{1}_g}} \text{RSP}^*$$



LEE-witnessed coinductive proof over $\text{Mil}^- + \{e = f \cdot e + g\}$

Proof transformation Mil \mapsto cMil

$$\begin{array}{c}
 \cdot \\
 \frac{e = f \cdot e + g}{e = f^* \cdot g} \text{RSP}^* \\
 \Longrightarrow \frac{e = f \cdot e + g \quad \mathcal{LCP}_{\text{Mil}^- + \{e = f \cdot e + g\}}(e = f^* \cdot g)}{e = f^* \cdot g} \text{LCoProof}_1
 \end{array}$$

Proof transformation $\text{Mil} \mapsto \text{cMil}$

Theorem

$\text{Mil} \approx \text{cMil}$, because:

every derivation in Mil with conclusion $e = f$ can be transformed effectively into a derivation in cMil with conclusion $e = f$.

Proof idea.

$$\frac{e = f \cdot e + g}{e = f^* \cdot g} \text{RSP}^* \quad \Longrightarrow \quad \frac{e = f \cdot e + g \quad \mathcal{LCP}_{\text{Mil}^- + \{e=f \cdot e + g\}}(e = f^* \cdot g)}{e = f^* \cdot g} \text{LCoProof}_1$$

Corollary

$\text{Mil} \sim \text{cMil} \sim \text{CLC}$.

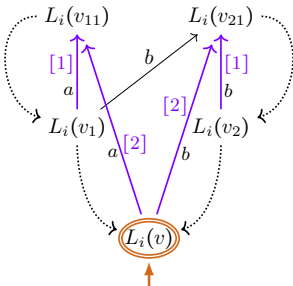
Proof transformation $\text{cMil} \mapsto \text{Mil}$

Lemma

For all star expression e, f , and equations $\Gamma \subseteq =_{\text{Mil}}$:

$$e \stackrel{\text{LEE}}{=}_{\text{Mil} + \Gamma} f \quad \Longrightarrow \quad e =_{\text{Mil}} f$$

Extraction of Mil-derivation from LEE-witn. coind. proof

 \underline{C}, \hat{C}


$$L_i(v_{21}) =_{\text{Mil}^-}^{(\text{sol})} 1 \cdot L_i(v_2) =_{\text{Mil}^-} L_i(v_2)$$

($=_{\text{Mil}^-}^{(\text{sol})}$ means use of 'is Mil^- -provable solution')

$$L_i(v_2) =_{\text{Mil}^-}^{(\text{sol})} b \cdot L_i(v_{21}) + 1 \cdot L_i(v) =_{\text{Mil}^-} b \cdot L_i(v_2) + L_i(v)$$

\Downarrow applying **RSP***

$$L_i(v_2) =_{\text{Mil}} b^* \cdot L_i(v)$$

$$L_i(v_{11}) =_{\text{Mil}^-}^{(\text{sol})} 1 \cdot L_i(v_1) =_{\text{Mil}^-} L_i(v_1)$$

$$\begin{aligned} L_i(v_1) &=_{\text{Mil}^-} a \cdot L_i(v_{11}) + b \cdot L_i(v_{21}) + 1 \cdot L_i(v) \\ &=_{\text{Mil}} a \cdot L_i(v_1) + (b \cdot b^* + 1) \cdot L_i(v) \\ &=_{\text{Mil}^-} a \cdot L_i(v_1) + b^* \cdot L_i(v) \end{aligned}$$

\Downarrow applying **RSP***

$$L_i(v_1) =_{\text{Mil}} a^* \cdot (b^* \cdot L_i(v)) =_{\text{Mil}^-} (a^* \cdot b^*) \cdot L_i(v)$$

$$\begin{aligned} L_i(v) &=_{\text{Mil}^-}^{(\text{sol})} 1 + a \cdot L_i(v_{11}) + b \cdot L_i(v_{21}) =_{\text{Mil}^-} 1 + a \cdot L_i(v_1) + b \cdot L_i(v_2) \\ &=_{\text{Mil}} (a \cdot (a^* \cdot b^*) + b \cdot b^*) \cdot L_i(v) + 1 \end{aligned}$$

\Downarrow applying **RSP***

$$L_i(v) =_{\text{Mil}} (a \cdot (a^* \cdot b^*) + b \cdot b^*)^* \cdot 1 =_{\text{Mil}^-} (a \cdot (a^* \cdot b^*) + b \cdot b^*)^* = \underline{s}_{\hat{C}}(v)$$

Proof transformation $\text{cMil} \mapsto \text{Mil}$

Lemma (extraction and unique solvability)

Let \underline{C} be a LEE-1-chart.

- ▶ From \underline{C} a Mil^- - (hence Mil -) provable solution can be extracted.
- ▶ Any two Mil -provable solutions of \underline{C} are Mil -provably equal.

Lemma

For all star expression e, f , and equations $\Gamma \subseteq =_{\text{Mil}}$:

$$e \stackrel{\text{LEE}}{=}_{\text{Mil}^+ + \Gamma} f \quad \Longrightarrow \quad e =_{\text{Mil}} f$$

Theorem

$\text{cMil} \lesssim \text{Mil}$, because:

every derivation in cMil with conclusion $e = f$ can be transformed effectively into a derivation in Mil with conclusion $e = f$.

Summary

We define:

- ▶ (LEE-witnessed) coinductive proofs over Mil^- :
 - ▶ 1-charts \underline{C} (with LEE) whose vertices are labeled by equations between the values of two provable solutions of \underline{C}
- ▶ proof systems
 - ▶ systems cMil / CLC with LEE-witnessed coind. proofs over Mil^-
 - ▶ systems \overline{cMil} / CC with coinductive proofs over Mil^-

Results:

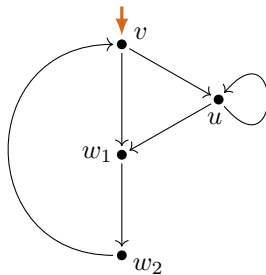
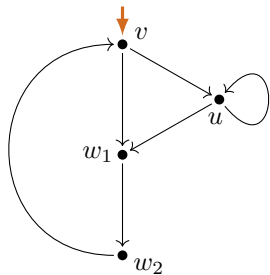
- ▶ $Mil = (Mil^- + RSP^*) \sim (Mil^- + LCoProof) = cMil \sim CLC$
- ▶ $Mil \approx (Mil^- + USP) \sim (Mil^- + CoProof) = \overline{cMil} \sim CC$ ((clearly) complete).

Desired application: proof strategy for completeness proof of Mil

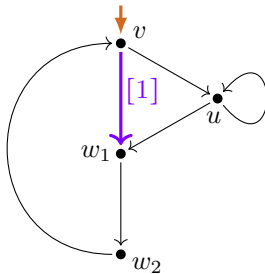
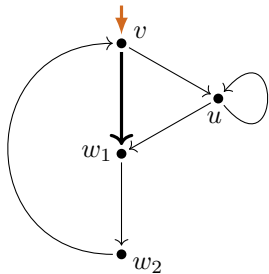
$$\vdash_{Mil} e = f \iff \vdash_{cMil} e = f \iff e = \llbracket \cdot \rrbracket_P f$$

- ▶ **Technical report:** [arXiv:2108.13104](https://arxiv.org/abs/2108.13104)

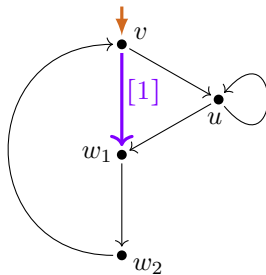
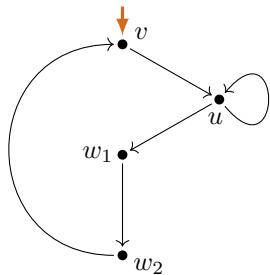
Layered LEE-witness



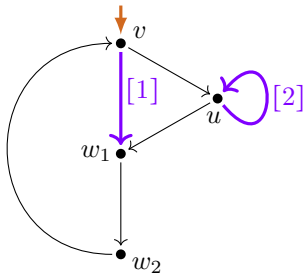
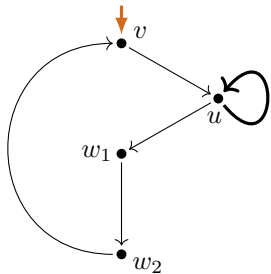
Layered LEE-witness



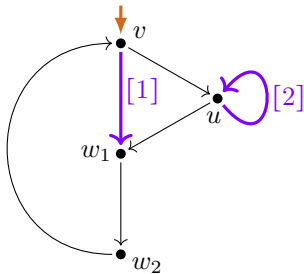
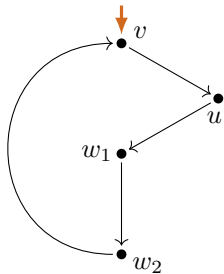
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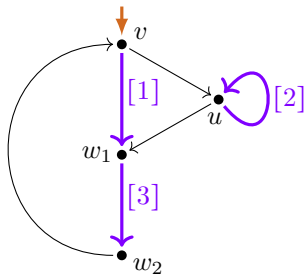
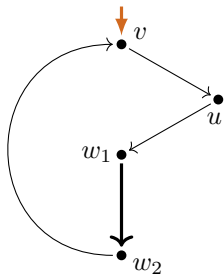
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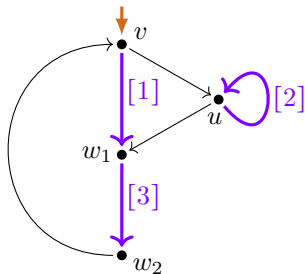
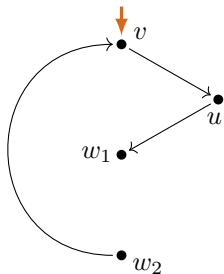
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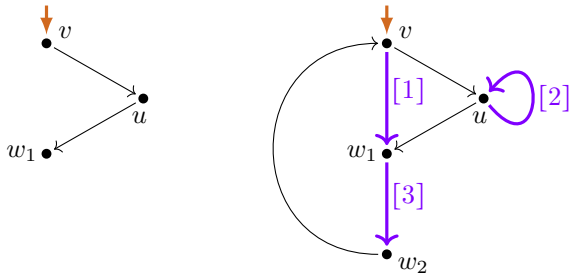
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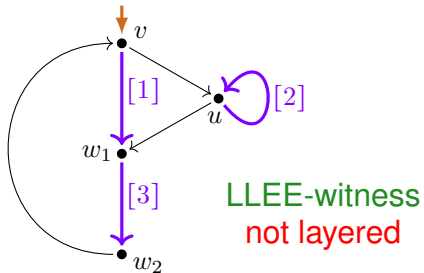
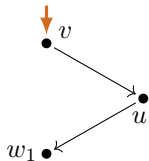
Layered LEE-witness



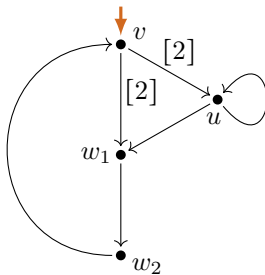
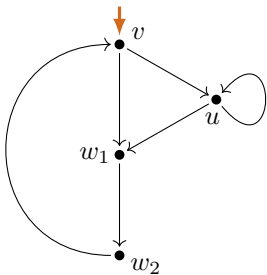
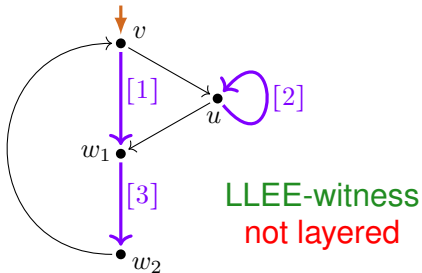
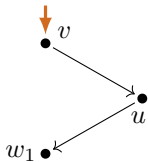
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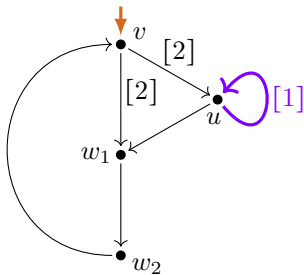
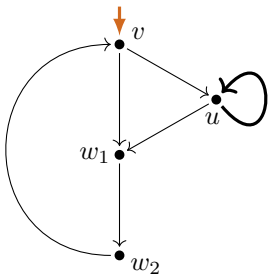
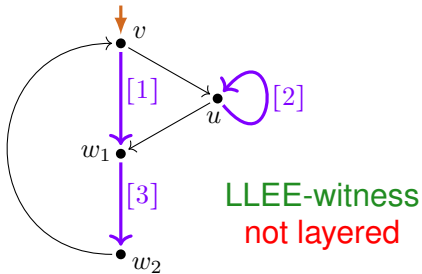
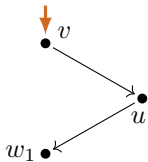
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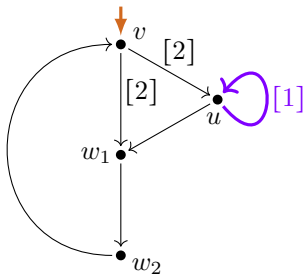
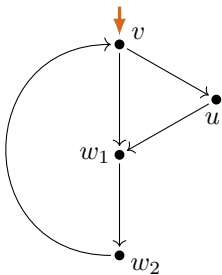
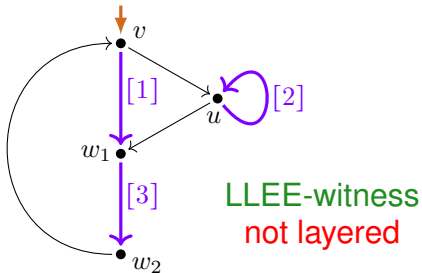
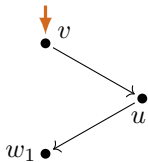
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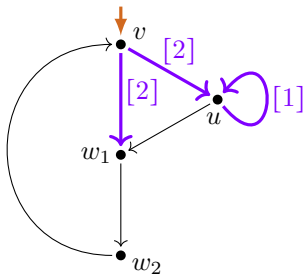
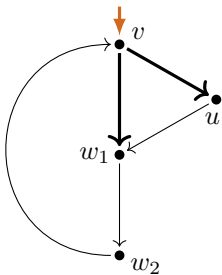
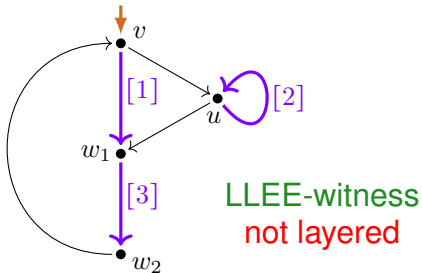
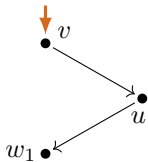
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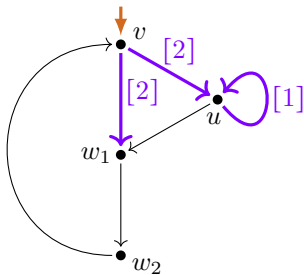
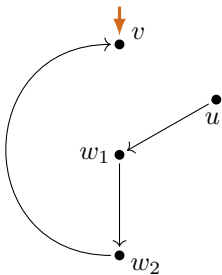
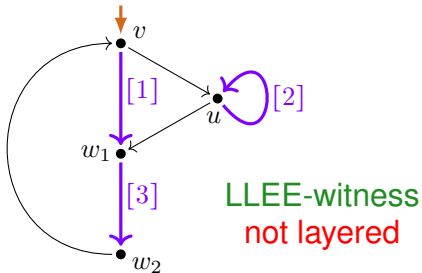
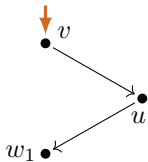
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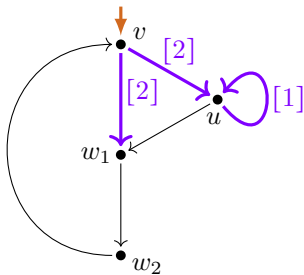
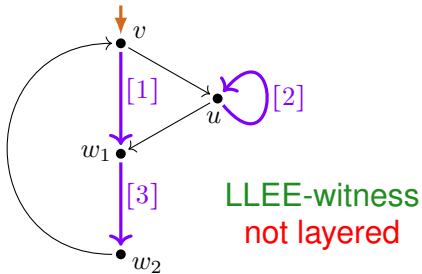
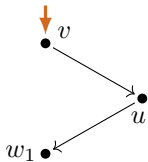
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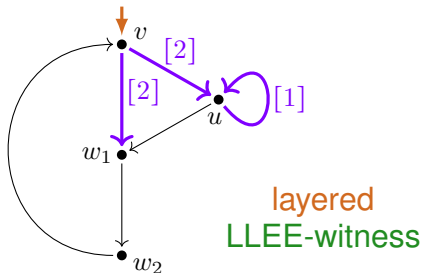
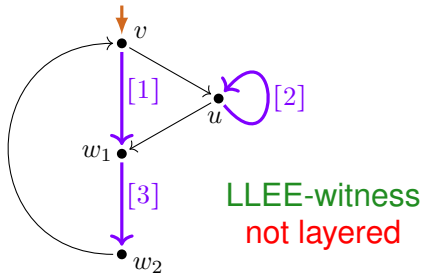
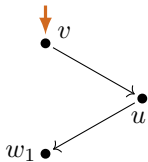
Layered LEE-witness



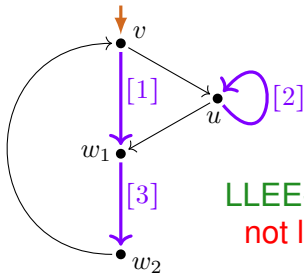
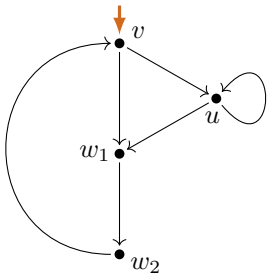
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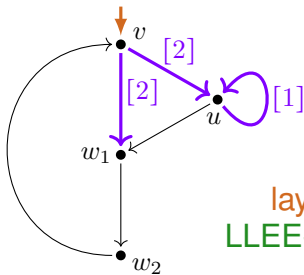
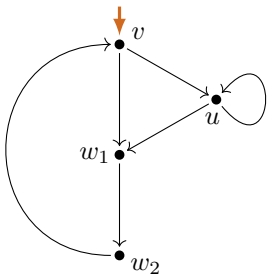
Layered LEE-witness



Layered LEE-witness

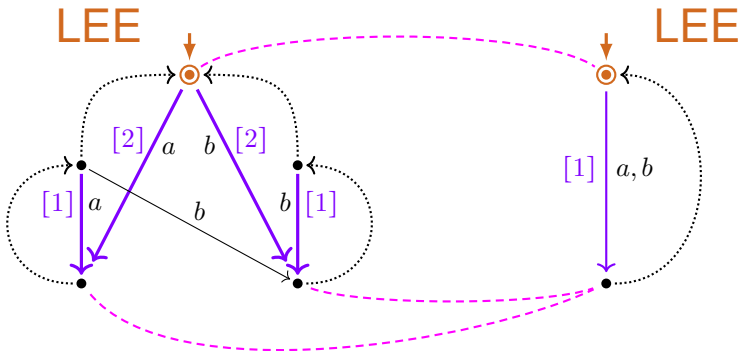


LLEE-witness
not layered



layered
LLEE-witness

LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

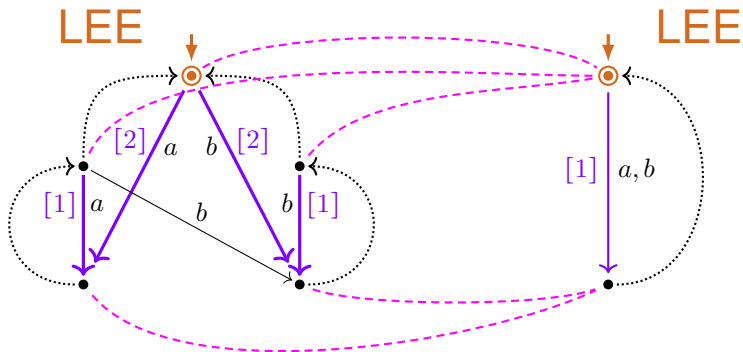


$$\underline{\mathcal{C}}((a^* \cdot b^*)^*)$$

$$\underline{\leftrightarrow}^{(1)}$$

$$\underline{\mathcal{C}}((a + b)^*)$$

LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

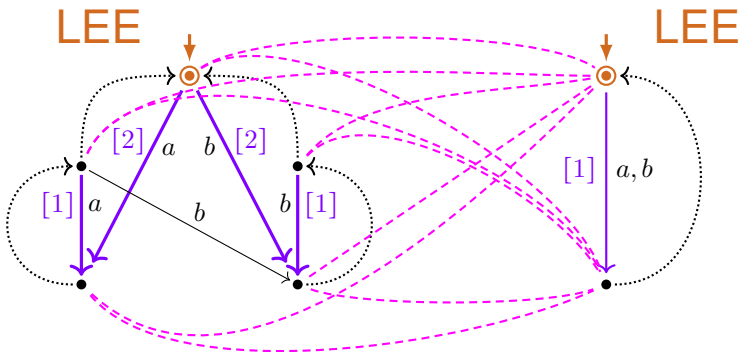


$$\underline{\mathcal{C}}((a^* \cdot b^*)^*)$$

$$\longleftrightarrow^{(1)}$$

$$\underline{\mathcal{C}}((a + b)^*)$$

LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

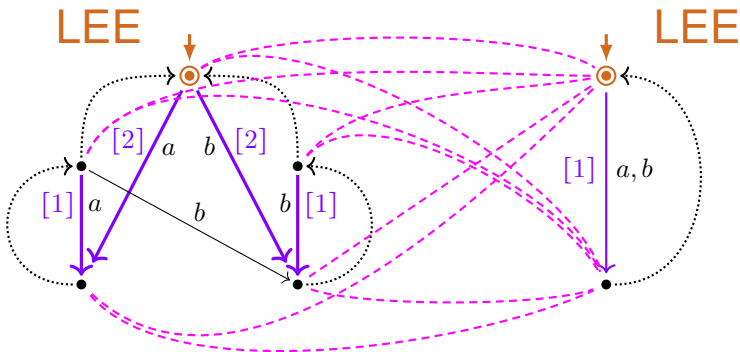


$$\underline{\mathcal{C}}((a^* \cdot b^*)^*)$$

$$\underline{\leftrightarrow}^{(1)}$$

$$\underline{\mathcal{C}}((a + b)^*)$$

LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

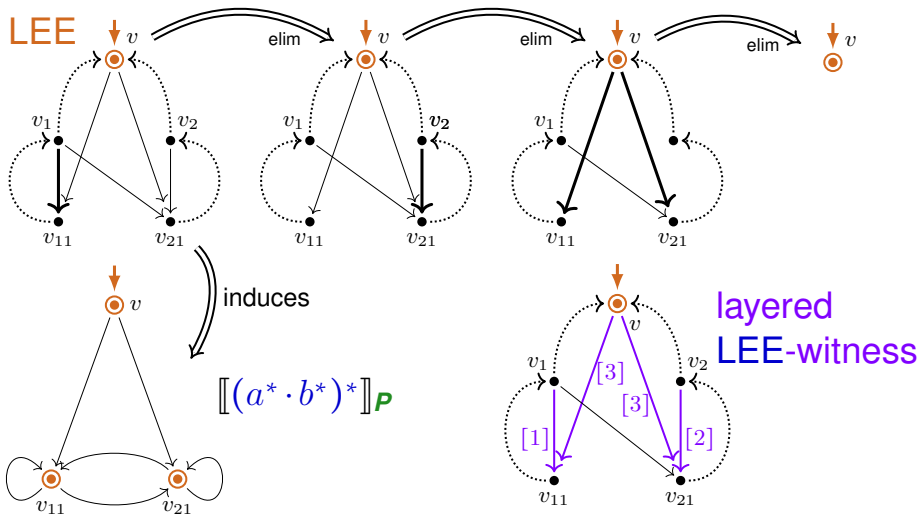


$$\underline{\mathcal{C}}((a^* \cdot b^*)^*)$$

$$\longleftrightarrow^{(1)}$$

$$\underline{\mathcal{C}}((a + b)^*)$$

LEE, and LLEE-witness, induced process graph



LEE-charts: properties and results

Lemmas

- (I) Chart interpretations of **1-free star expressions** satisfy **LEE**.
- (SU) **LEE**-charts have unique provable solutions up to **Mil**-provability.
- (C) **LEE** is preserved under bisimulation collapse.

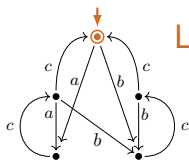
Theorem (G/Fokkink, LICS 2020)

The adaptation **BBP** of **Mil** to **1-free star expressions** is complete.

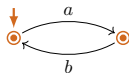
Consequence of lemmas used

- (E) A chart \mathcal{C} is **expressible by a 1-free star expr. modulo bisimilarity**
 \iff the bisimulation collapse of \mathcal{C} is a **LEE**-chart.

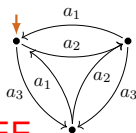
Hence expressible | not expressible by **1-free star expressions**:



LEE



LEE



LEE